

THE OPERATIONAL PERSPECTIVE

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Advances in Proof Theory
In honor of Gerhard Jäger's 60th birthday
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Operationally Based Axiomatic Programs

- The Explicit Mathematics Program
- The Unfolding Program
- A Logic for Mathematical Practice
- Operational Set Theory (OST)

Foundations of Explicit Mathematics

- Book in progress with Gerhard Jäger and Thomas Strahm, with the assistance of Ulrik Buchholtz
- An online bibliography

The Unfolding Program

- Open-ended Axiomatic Schemata; language not fixed in advance
- Examples in Logic, Arithmetic, Analysis, Set Theory
- The general concept of unfolding explained within an operational framework

Aim of the Unfolding Program

- S an open-ended schematic axiom system
- Which operations on individuals--and which on predicates--and what principles concerning them ought to be accepted once one has accepted the operations and principles of S?

Results on (Full) Unfolding

- Non-Finitist Arithmetic (NFA);
 $|\mathcal{U}(\text{NFA})| = \Gamma_0$
- Finitist Arithmetic (FA):
 $\mathcal{U}(\text{FA}) \equiv \text{PRA}, \mathcal{U}(\text{FA} + \text{BR}) \equiv \text{PA}$
- (Feferman and Strahm 2000, 2010)

Unfolding of ID_1

- $|U(ID_1)| = \psi(\Gamma_{\Omega+1})$
(U. Buchholtz 2013)
- Note: $\psi(\Gamma_{\Omega+1})$ is to $\psi(\varepsilon_{\Omega+1})$ as Γ_0 is to ε_0 .

Problems for Unfolding to Pursue

- Unfolding of analysis
- Unfolding of KP + Pow
- Unfolding of set theory

Indescribable Cardinals and Admissible Analogues Revisited

- Aim: To have a straightforward and principled transfer of the notions of indescribable cardinals from set theory to admissible ordinals.
- A new proposal and several conjectures, suggested at the end of the OST paper.
- NB: Not *within* OST

Aczel and Richter Pioneering Work

- Aczel and Richter [A-R] (1972)
Richter and Aczel [R-A] (1974)
- In set theory, assume κ regular $> \omega$.
- Let $f, g: \kappa \rightarrow \kappa$; $F(f) = g$ type 2 over κ .

[A, R]-2

- F is *bounded* $\Leftrightarrow (\forall f: \kappa \rightarrow \kappa)(\forall \xi < \kappa)$
[$F(f)(\xi)$ is det. by $< \kappa$ values of f]
- α is a *witness* for $F \Leftrightarrow (\forall f: \kappa \rightarrow \kappa)$
[$f: \alpha \rightarrow \alpha \Rightarrow F(f): \alpha \rightarrow \alpha$]
- κ is *2-regular* iff every bounded F has a witness.

[A, R]-3

- Notions of bounded, witness, n -regular for $n > 2$ are “defined in a similar spirit”, but never published.
- Theorem 1. K is $n+1$ -regular iff K is strongly Π^1_n -indescribable.
- Proved only for $n = 1$ in [R-A](1974).

[A, R]-4

- Admissible analogues:
- Assume κ admissible $> \omega$
- κ is *n-admissible*, obtained by replacing ‘bounded’ in the defn. of n-regular by ‘recursive’, functions by their Gödel indices, and functionals by recursive functions applied to such indices.

[A, R]-5

- Theorem 2. κ is n -admissible iff κ is Π^0_{n+1} reflecting.
- Proved only for $n = 2$ in [R-A](1974).
- Proposed:
Least Π^0_{n+2} -reflecting ordinal \sim least [strongly] Π^1_n -indescribable cardinal.

A Proposed New Approach

- Directly lift to card's and admissible ord's notions of *continuous functionals of finite type* from o.r.t.
- Kleene (1959), Kreisel (1959)
- Deal only with objects of *pure type n*.
- $\mathcal{K}^{(0)} = \mathcal{K}$; $\mathcal{K}^{(n+1)} = \text{all } F^{(n+1)}: \mathcal{K}^{(n)} \rightarrow \mathcal{K}$.

“Sequence Numbers” in Set Theory

- Assume κ a strongly inaccessible cardinal.
- Let $\kappa^{<\kappa} =$ all sequences $s: \alpha \rightarrow \kappa$ for arbitrary $\alpha < \kappa$.
- Fix $\pi: \kappa^{<\kappa} \rightarrow \kappa$, one-one and onto; so $\pi(g \upharpoonright \alpha)$ is an ordinal that codes $g \upharpoonright \alpha$.

Continuous Functionals and Their Associates

- Inductive definition of $F \in C^{(n)}$, and of f is an associate of F , where f is of type 1:
- For $n = 1$, f is an associate of F iff $f = F$.
- For $F \in K^{(n+1)}$, f is an associate of F iff for every G in $C^{(n)}$ and every associate g of G ,

Continuous Functionals and Their Associates (cont'd)

- (i) $(\exists \alpha, \beta < \kappa)(\forall \gamma)[\alpha \leq \gamma < \kappa \Rightarrow f(\pi(g \upharpoonright \gamma)) = \beta + 1]$, and
- (ii) $(\forall \gamma, \beta < \kappa) [f(\pi(g \upharpoonright \gamma)) = \beta + 1 \Rightarrow F(G) = \beta]$.
- F is in $C^{(n+1)}$ iff F has some associate f .

Witnesses

- For F in $C^{(n)}$ and $\alpha < \kappa$, define α is a witness for F , as follows:
- For $n = 1$, and $F = f$, α is a witness for F iff $f : \alpha \rightarrow \alpha$.
- For $F \in C^{(n+1)}$, α is a witness for F iff $(\forall G \in C^{(n)}) [\alpha \text{ a witness for } G \Rightarrow F(G) < \alpha]$.

$C^{(n)}$ -Regularity; Conjectures

- κ is $C^{(n)}$ -reg for $n > 1$ iff every F in $C^{(n)}$ has some witness $\alpha < \kappa$.
- Conjecture 1. For each $n \geq 1$, the predicate *f is an associate of some F in $C^{(n+1)}$* , is definable in Π^1_n form.
- Conjecture 2. For each $n \geq 1$, κ is $C^{(n+1)}$ -reg iff κ is strongly Π^1_n -indescribable.
- Conj-2 holds for $n = 1$ by [R-A] proof.

Analogues over Admissibles

- Consider admissible $\kappa > \omega$.
- For analogues in (κ -) recursion theory replace functions of type 1 by indices ζ of (total) recursive functions $\{\zeta\}$.
- But then at type 2 (and higher) we must restrict to those functions $\{\zeta\}$ that act extensionally on indices.

Effective Operations over Admissibles

- Following Kreisel (1959), define the class E_n of (κ -) *effective operations of type n* , and the relation \equiv_n by induction on $n > 0$:
- E_1 consists of all indices ζ of recursive functions;
 $\zeta \equiv_1 \nu$ iff for all ξ , $\{\zeta\}(\xi) = \{\nu\}(\xi)$.

Effective Operations over Admissibles (cont'd)

- $\zeta \in E_{n+1} \Leftrightarrow \{\zeta\}: E_n \rightarrow K$ and
 $(\forall \xi, \eta \in E_n)[\xi \equiv_n \eta \Rightarrow \{\zeta\}(\xi) = \{\zeta\}(\eta)]$;
 $\zeta \equiv_{n+1} v \Leftrightarrow (\forall \xi \in E_n)[\{\zeta\}(\xi) = \{v\}(\xi)]$.
- Conjecture 3. Every type $n+1$ effective operation is the restriction of a functional in $C^{(n+1)}$.
- This would show why can drop the boundedness hypothesis in analogue.

Witnesses for Effective Operations

- For ζ in E_1 , α is a witness for ζ iff $\{\zeta\}: \alpha \rightarrow \alpha$.
- For ζ in E_{n+1} when $n \geq 1$,
 α is a witness for $\zeta \Leftrightarrow (\forall \xi \in E_n)$
 $[\alpha \text{ a witness for } \xi \Rightarrow \{\zeta\}(\xi) < \alpha]$.
- κ is E_n -admissible if each ζ in E_n has some witness $\alpha < \kappa$. (Equiv. to $[A, R]$ n -admiss.)

Further Work

- Settle the conjectures.
- (Scott) The partial equivalence relation approach to types in λ -calculus models over $\mathcal{P}(\mathbb{N})$ gives a "clean" definition of the Kleene-Kreisel hierarchy. Can this idea be generalized to $\mathcal{P}(\kappa)$? [What about effective operations?]

Further Work (cont'd)

- The present approach leaves open the question as to what is the proper analogue for admissible ordinals--if any--of a cardinal κ being Π^m_n -indescribable for $m > 1$.

The End